

Optimizing **Abstract Abstract Machines**

J. Ian Johnson,

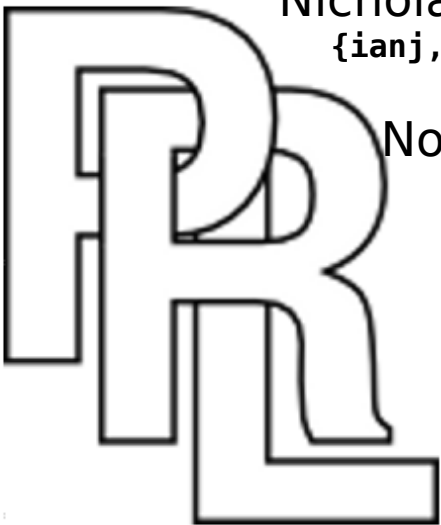
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Abstract
abstract machines?

Abstracting Abstract Machines

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Abstract

We describe a derivational approach to abstract interpretation that yields novel and transparently sound static analyses when applied to well-established abstract machines. To demonstrate the technique and support our claim, we transform the CEK machine of Felleisen and Friedman, a lazy variant of Krivine's machine, and the stack-inspecting CM machine of Clements and Felleisen into abstract interpretations of themselves. The resulting analyses bound temporal ordering of program events; predict return-flow and stack-inspection behavior; and approximate the flow and evaluation of by-need parameters. For all of these machines, we find that a series of well-known concrete machine refactorings, plus a technique we call store-allocated continuations, leads to machines that abstract into static analyses simply by bounding their stores. We demonstrate that the technique scales up uniformly to allow static analysis of realistic language features, including tail calls, conditionals, side effects, exceptions, first-class continuations, and even garbage collection.

Categories and Subject Descriptors F.3.2 [Logics and Meanings of Programs]: Semantics of Programming Languages—Program analysis, Operational semantics; F.4.1 [Mathematical Logic and Formal Languages]: Mathematical Logic—Lambda calculus and related systems

General Terms Languages, Theory

Keywords abstract machines, abstract interpretation

1. Introduction

Abstract machines such as the CEK machine and Krivine's machine are first-order state transition systems that represent the core of a real language implementation. Semantics-based program analysis, on the other hand, is concerned with safely approximating intensional properties of such a machine as it runs a program. It seems natural then to want to systematically derive analyses from machines to approximate the core of realistic run-time systems.

Our goal is to develop a technique that enables direct abstract interpretations of abstract machines by methods for transforming a given machine description into another that computes its finite approximation.

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We demonstrate that the technique of refactoring a machine with **store-allocated continuations** allows a direct structural abstraction¹ by bounding the machine's store. Thus, we are able to convert semantic techniques used to model language features into static analysis techniques for reasoning about the behavior of those very same features. By abstracting well-known machines, our technique delivers static analyzers that can reason about by-need evaluation, higher-order functions, tail calls, side effects, stack structure, exceptions and first-class continuations.

The basic idea behind store-allocated continuations is not new. SML/NJ has allocated continuations in the heap for well over a decade [28]. At first glance, modeling the program stack in an abstract machine with store-allocated continuations would not seem to provide any real benefit. Indeed, for the purpose of defining the meaning of a program, there is no benefit, because the meaning of the program does not depend on the stack-implementation strategy. Yet, a closer inspection finds that store-allocating continuations eliminate recursion from the definition of the state-space of the machine. With no recursive structure in the state-space, an abstract machine becomes eligible for conversion into an abstract interpreter through a simple structural abstraction.

To demonstrate the applicability of the approach, we derive abstract interpreters of:

- a call-by-value λ -calculus with state and control based on the CESK machine of Felleisen and Friedman [13],
- a call-by-need λ -calculus based on a tail-recursive, lazy variant of Krivine's machine derived by Ager, Danvy and Midtgaard [1], and
- a call-by-value λ -calculus with stack inspection based on the CM machine of Clements and Felleisen [3];

and use abstract garbage collection to improve precision [25].

Overview

In Section 2, we begin with the CEK machine and attempt a structural abstract interpretation, but find ourselves blocked by two recursive structures in the machine: environments and continuations. We make three refactorings to:

1. store-allocate bindings,
2. store-allocate continuations, and
3. time-stamp machine states;

resulting in the CESK, CESK*, and time-stamped CESK* machines, respectively. The time-stamps encode the history (context) of the machine's execution and facilitate context-sensitive abstractions. We then demonstrate that the time-stamped machine abstracts directly into a parameterized, sound and computable static analysis.

¹ A structural abstraction distributes component-, point-, and member-wise.

Abstracting Abstract Machines

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that a series of well-known techniques we call store-allocated continuations that abstract into static analyses. We demonstrate that static analysis of recursive conditionals, side effects, even garbage collection

Categories and Subject of Programs: Static analysis, Operational Formal Languages, related systems

General Terms: Languages

Keywords: abstract machines

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program stack in an abstract machine. Such abstractions would not seem to require the re-derivation of defining the meaning of the implementation strategy for allocating continuations of the state-space of the state-space, an abstraction into an abstract machine.

In this approach, we derive a control based on the [13].

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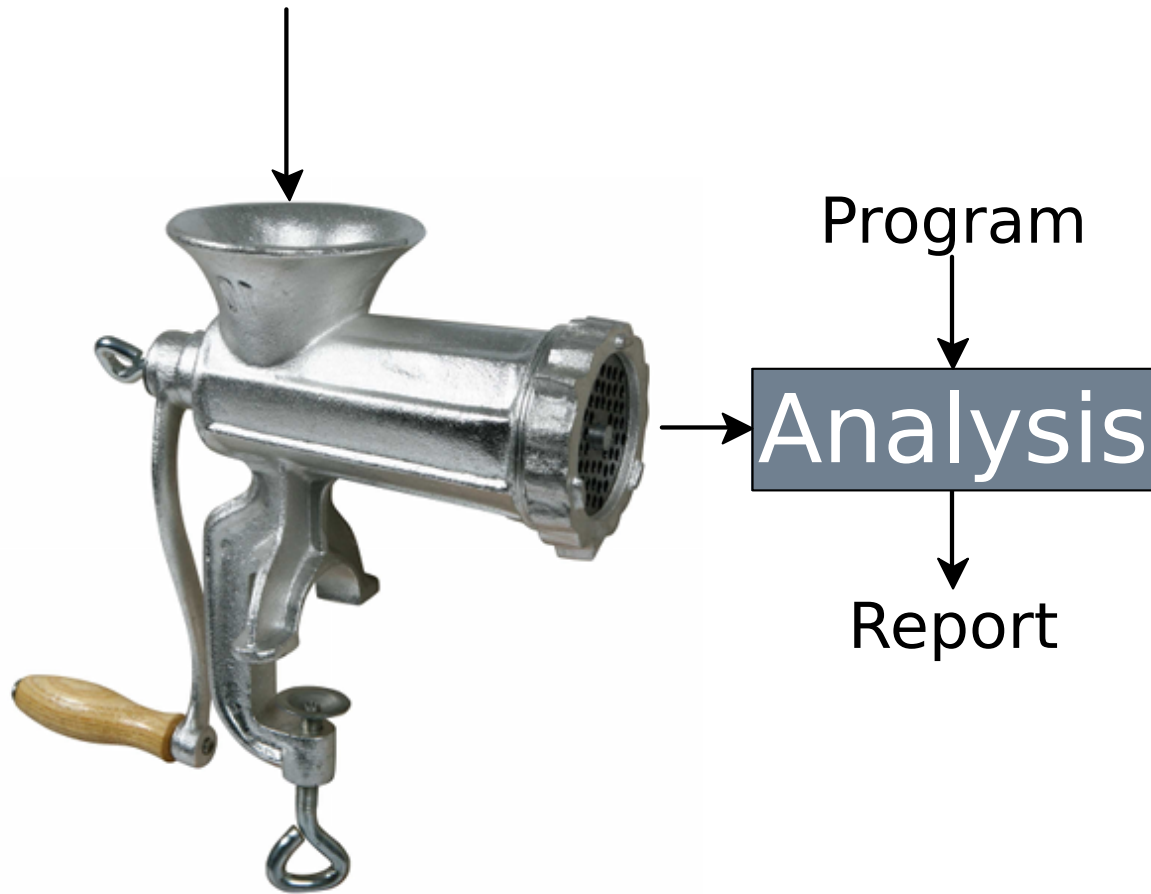
s?

Semantics



→ Analysis

Semantics



Semantics



Program

Analysis

Maybe

Goal: Directly implementable math

Competitive performance,*
minimal effort,
simple proofs

same ballpark, nosebleed seats

Goal: Directly implementable math

Competitive performance,*

Slogan:

Engineering tricks
as semantics refactorings

same ballpark, nosebleed seats

An ideal static analysis has

Performance

An ideal static analysis has

Soundness

Maintainability

Precision

Performance

An ideal static analysis has

[Van Horn and Might ICFP 2010]
(AAM)

[Soundness
Maintainability
Precision
Performance

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[Van Horn and Might ICFP 2010]
(AAM)

[Soundness
Maintainability
Precision
Performance]

[Johnson, et al. ICFP 2013]
(OAAM)

OAAM outline

OAAM outline

- Store-allocate values
- Frontier-based semantics
- Lazy non-determinism
- Abstract compilation
- Locally log-based store-deltas
- Store-counting
- Mutable frontier and store

OAAM outline

Decrease state space

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Explore faster

OAAM outline

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Explore faster

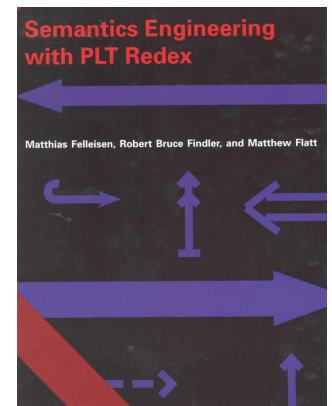
First: recap of AAM

Start with CESK machine

Control



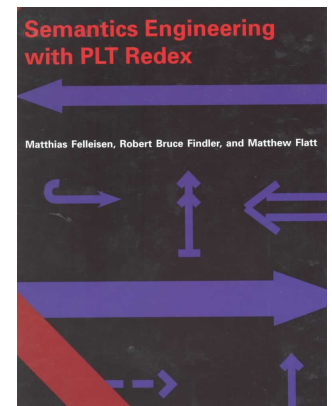
$\langle e, \rho, \sigma, \kappa \rangle$



Start with CESK machine

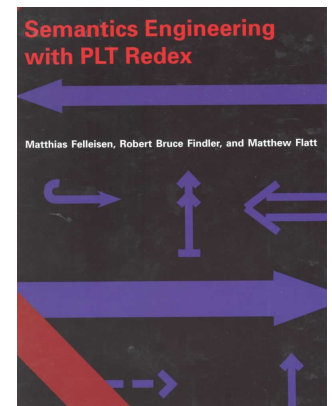
$$\langle e, \rho, \sigma, \kappa \rangle$$

Environment



Start with CESK machine

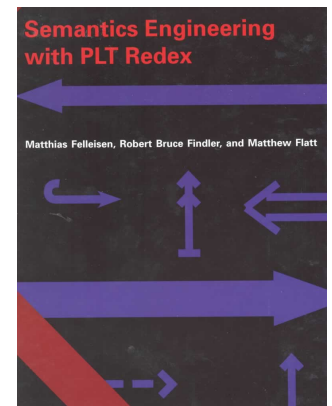
Store
↓
 $\langle e, \rho, \sigma, \kappa \rangle$



Start with CESK machine

$\langle e, \rho, \sigma, \kappa \rangle$

Kontinuation



$$\rho \in \text{Env} \quad = \quad \text{Var} \rightarrow \text{Addr}$$

$$\kappa \in \text{Kont} \quad ::= \quad [] \mid \varphi : \kappa$$

$$\sigma \in \text{Store} \quad = \quad \text{Addr} \rightarrow (\text{Value} \times \text{Env})$$

$$\langle x, \rho, \sigma, \kappa \rangle \mapsto \langle v, \rho', \sigma, \kappa \rangle$$

$$\text{if } (v, \rho') = \sigma(\rho(x))$$

$$\langle (e_0 \ e_1), \rho, \sigma, \kappa \rangle \mapsto \langle e_0, \rho, \sigma, \text{ar}(e_1, \rho) : \kappa \rangle$$

$$\langle v, \rho, \sigma, \text{ar}(e, \rho) : \kappa \rangle \mapsto \langle e, \rho, \sigma, \text{fn}(v, \rho) : \kappa \rangle$$

$$\langle v, \rho, \sigma, \text{fn}(\lambda x. e, \rho') : \kappa \rangle \mapsto \langle e, \rho'', \sigma', \kappa \rangle$$

$$\text{where } \rho'' = \rho' [x \mapsto a]$$

$$\sigma' = \sigma [a \mapsto (v, \rho)]$$

a fresh

$$\rho \in \text{Env} = \text{Var} \rightarrow \text{Addr}$$

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$$\langle (e_0 e_1), \rho, \sigma, \kappa \rangle \mapsto \langle e_0, \rho, \sigma', \text{ar}(e_1, \rho) : a \rangle \quad \sigma' = \sigma[a \mapsto \kappa]$$

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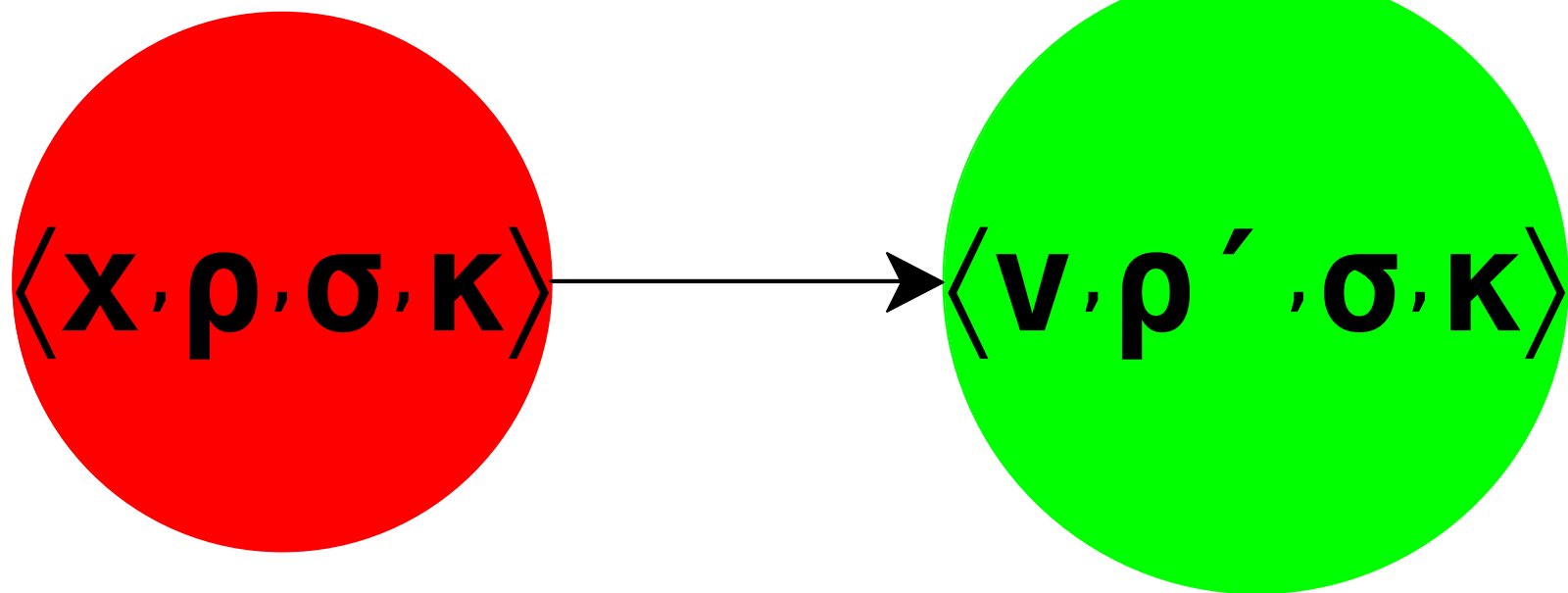
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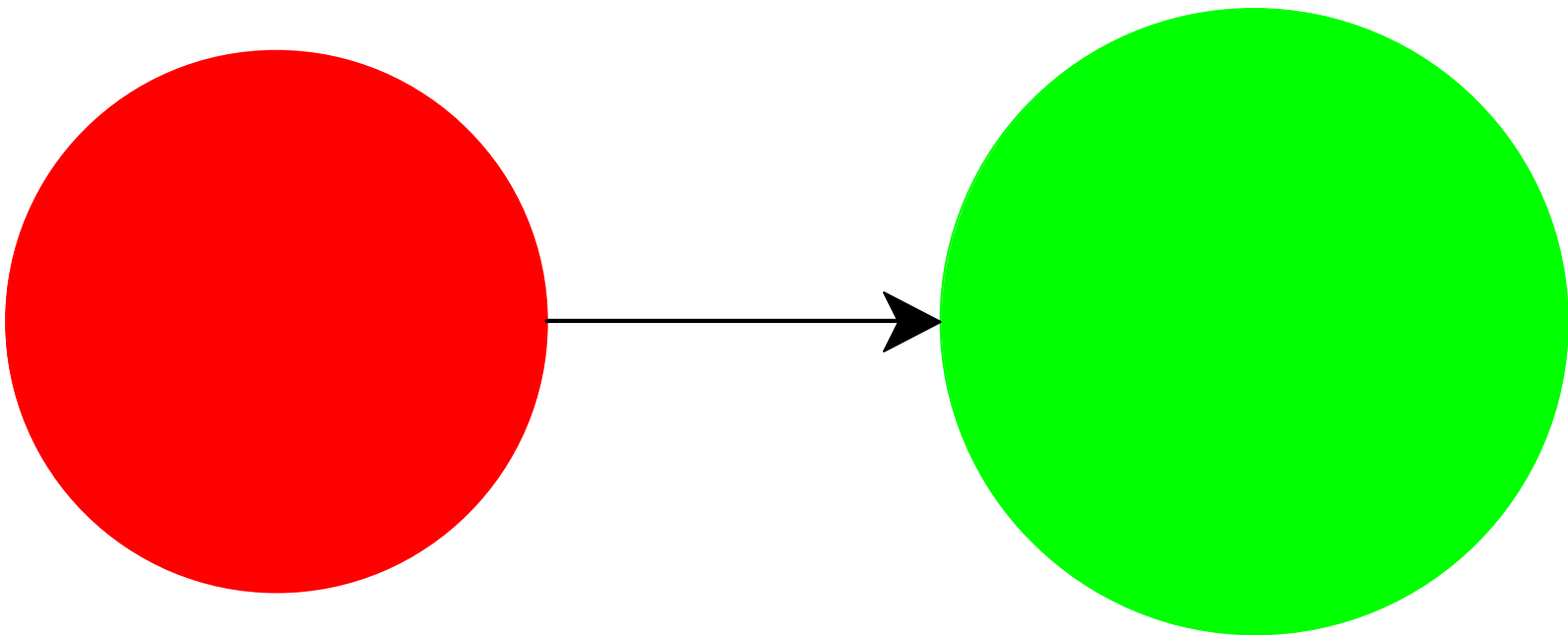
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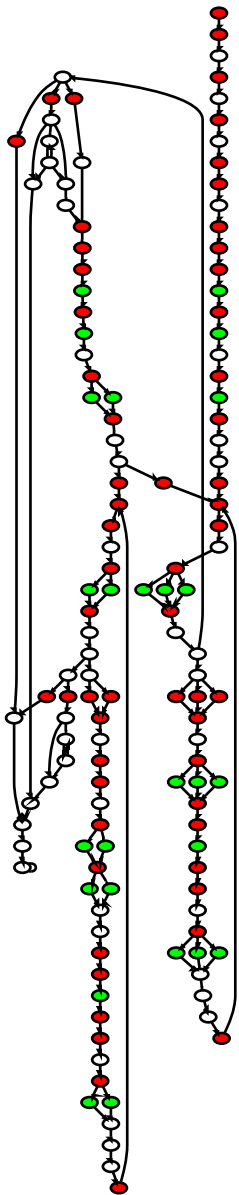
$$a = \text{alloc}(\varsigma)$$

$$\langle \mathbf{x}, \boldsymbol{\rho}, \sigma, \mathbf{K} \rangle \mapsto \langle \mathbf{v}, \boldsymbol{\rho}', \sigma, \mathbf{K} \rangle$$

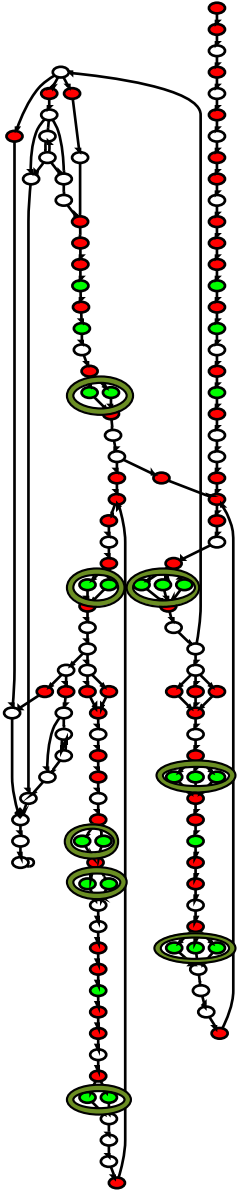




Finite reduction relation = graph

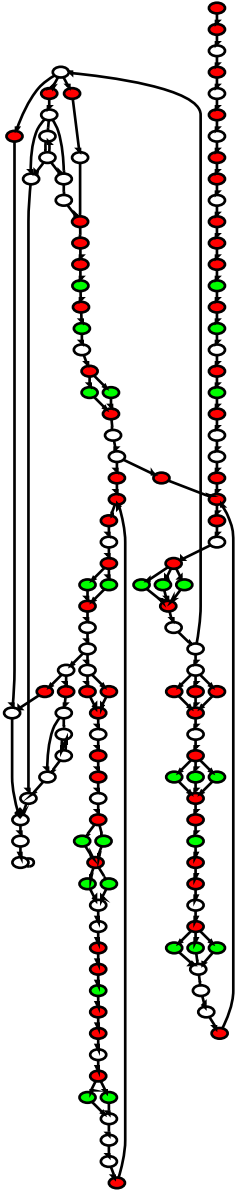


Finite reduction relation = graph

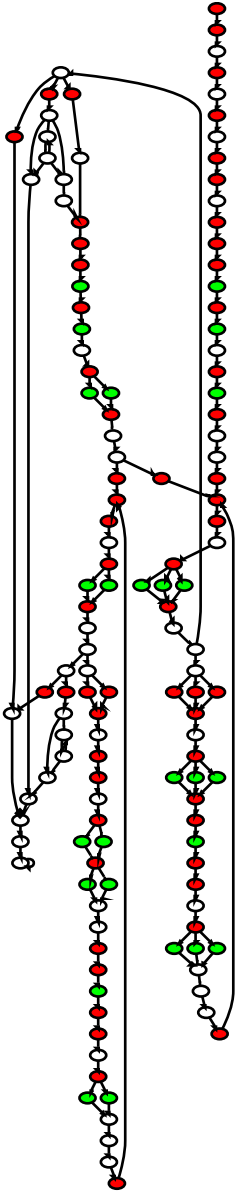


1. Lazy non-determinism

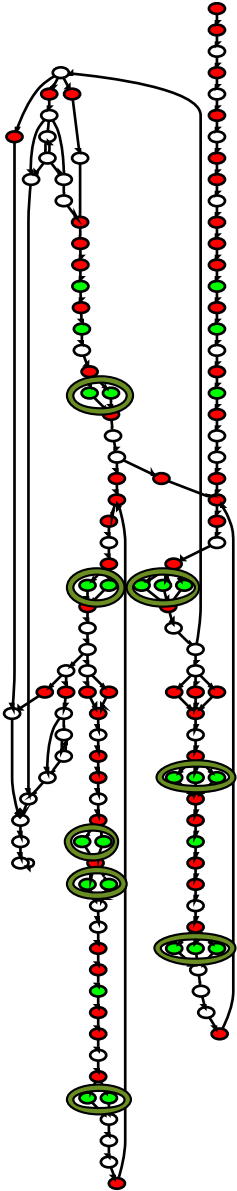
Finite reduction relation = graph



1. Lazy non-determinism
2. Abstract compilation



1. Lazy non-determinism
2. Abstract compilation
3. Store deltas



1. Lazy non-determinism
2. Abstract compilation
3. Store deltas

$$\rho \in \text{Env} = \text{Var} \rightarrow \text{Addr}$$

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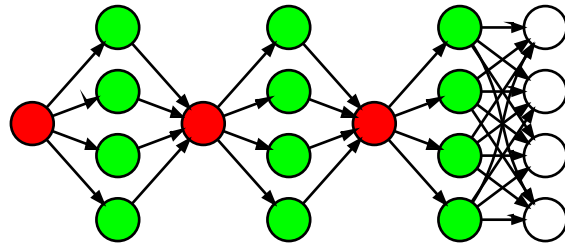
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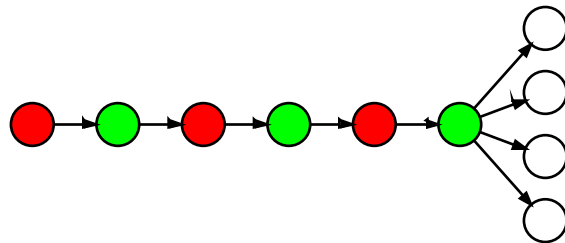
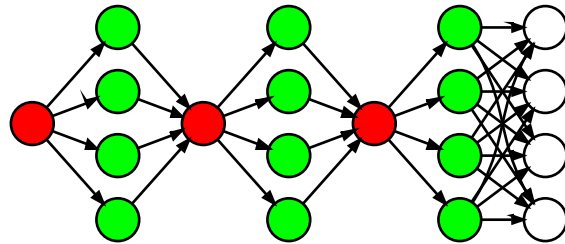
$$\sigma' = \sigma \sqcup [a \mapsto \text{force}(v)]$$

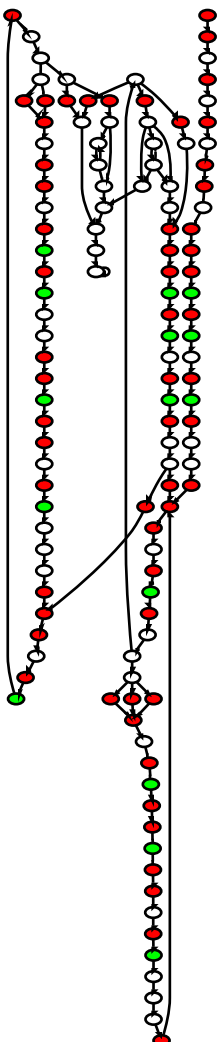
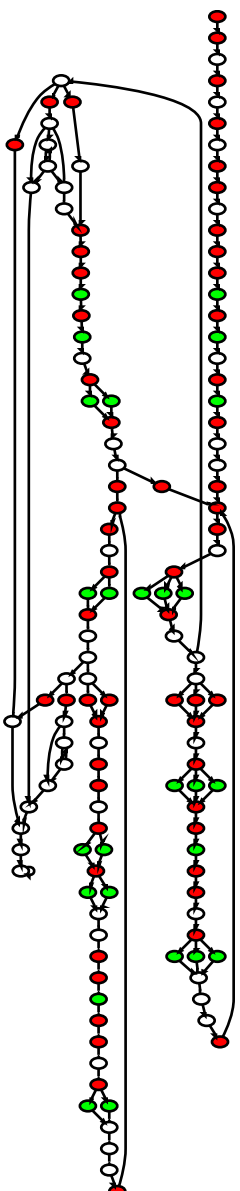
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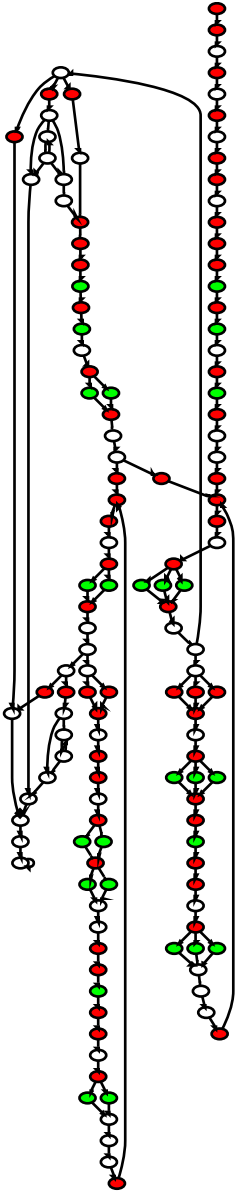
$(f \ x \ y)$



$(f \times y)$







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```
(match e ((var x) ((continue v σ κ) : v ∈ (get σ (get ρ x))))) ((
```

```
(match x (ifw x) (continue x w) : x < (get x (get p (lisp w x)) (eval w p + (cons (w x p) x))) (lisp x w*) (continue (closure x p) w x)) (use) (Anin placet Cosby beer party on a Anin readmade American placet brunch Cupidatet beard put))
(Cupidatet it on sweater pour-over x and Cosby (Anin placet sweater on a placet beard on Apparel rights DIY put American doloer selvaage on Parlatier odia on Leggings American it Cupidatet use))
(Cupidatet Truffaut) (pour-over Truffaut dreamcatcher Anin bird cupidatet party it Brooklyn American it bird placet Truffaut brunch et Truffaut denim brunch booth bicycle DIY bird on it artisan readmade American))
(bird roof doloer a sud put brunch et beard on) (leggings it Truffaut Parlatier Cosby raw on dreamcatcher palen roof party rights pour-over leggings placemat party Truffaut Ania palen odia getrify readmade cupidatet placemat pour-over doloer fingerstache Ania bird))
(Cupidatet artisan) (fingerstache Craft pour-over odia on cupidatet bird Truffaut Cosby put bird a sud photo leggings sweater Apparel rights a fingerstache Craft Cosby and Leggings American))
(Brunch bird) (placemat palen a on put a booth and DIY getrify American Aliqua Parlatier brunch odia placemat))
(Parlatier Excepteur doloer Aliqua pour-over sud bird a it) (roof Craft fingerstache booth put on Truffaut on placemat Parlatier DIY getrify party dreamcatcher Parlatier put sud Laborum getrify put))
(Parlatier bird it Excepteur Apparel booth bicycle) (brunch))
(denim put sud leggings fingerstache) (party placemat raw cupidatet Brooklyn beard roof party getrify placemat raw sud artisan beard Excepteur leggings on Cosby American denim DIY Truffaut booth Laborum roof Aliqua beer Apparel put))
(bicycle booth American) (et put Laborum et Aliqua getrify raw readmade Parlatier))
(heard raw odia a it fingerstache) (sweater Apparel pour-over leggings roof Cupidatet Craft rights Parlatier Truffaut Parlatier doloer bird it odia Apparel dreamcatcher Ania it it et party denim))
(Apparel bicycle selvaage bird readmade DIY DIY a bicycle) (put Parlatier put getrify leggings brunch sud on Parlatier odia Brooklyn bicycle Aliqua et on odia))
(denim Craft sud brunch odia) (Excepteur selvaage Parlatier sweater beer bicycle roof pour-over readmade photo Brooklyn Parlatier dreamcatcher Aliqua Brooklyn on leggings party cupidatet party denim DIY Excepteur it denim doloer Cupidatet dreamcatcher Ania))
(it readmade DIY rights artisan leggings party raw party put) (liscant it on booth beer rights it Apparel bicycle getrify a beard Craft rights and booth put selvaage Laborum American Cosby Laborum artisan pour-over sweater))
(sud pour-over readmade doloer pour-over it) (Aliqua selvaage DIY on bicycle booth Truffaut sud and roof Cosby fingerstache it pour-over on getrify dreamcatcher artisan denim et bicycle bird dreamcatcher roof leggings))
(placemat Laborum) (a Truffaut roof raw getrify Parlatier denim party a placemat a sweater))
(sweater it leggings) (elssomed fingerstache put Cupidatet a American dreamcatcher leggings Parlatier placemat bicycle bird artisan Craft a American readmade beard placemat on))
(dolor sweater American selvaage beard on artisan placemat) (DIY))
(Laborum getrify odia Apparel a booth leggings leggings) (Cupidatet bicycle cupidatet elssomed fingerstache it doloer dreamcatcher raw roof bird artisan Aliqua and DIY odia it Aliqua American elssomed Apparel leggings artisan bird it it elssomed))
(party Apparel cupidatet brunch) (beer sud on Aliqua placemat palen readmade Parlatier raw Apparel Apparel sweater leggings Ania rights on on roof Cosby readmade sud on booth Apparel it Craft Brooklyn odia))
(sweater American artisan party a raw sud pour-over sud) (bicycle Laborum beard beard dreamcatcher))
(roof) (bicycle party cupidatet DIY denim artisan Parlatier Cupidatet Cupidatet and Truffaut beard fingerstache readmade cupidatet elssomed beard it leggings American bicycle fingerstache DIY DIY leggings sud))
(American party Parlatier odia sud rights and bird Craft) (roof denim selvaage pour-over on brunch sweater Brooklyn bird it brunch photo))
(Laborum) (artisan odia Aliqua American a cupidatet bird put artisan it palen palen Apparel Apparel Ania it Apparel rights Cupidatet bird doloer a doloer doloer Ania))
(it bicycle bicycle Ania palen placemat) (Aliqua Brooklyn Excepteur a Cupidatet))
(photo) (sweater bird sud it DIY palen it leggings leggings Apparel sud bird American Truffaut pour-over Aliqua))
(denim artisan on doloer photo put et beard DIY beer) (Craft on it Cupidatet Excepteur getrify DIY leggings sweater beer bird put palen dreamcatcher beer Brooklyn Aliqua Brooklyn bird palen Parlatier Aliqua bird))
(Cosby brunch Laborum selvaage sud photo beer odia) (Brooklyn beard Aliqua denim Truffaut pour-over Aliqua et placemat sud Craft odia Cosby Apparel it Truffaut put sweater it bird dreamcatcher bird Truffaut it denim cupidatet))
(Cosby selvaage Parlatier it getrify getrify photo) (readmade denim sud put DIY sud readmade Laborum Cupidatet raw Laborum denim booth raw party Apparel photo Craft booth dreamcatcher bird cupidatet denim sud put))
(Brooklyn placemat a cupidatet palen Cosby selvaage doloer) (bird it on beer on bird booth Excepteur Apparel a fingerstache put artisan Parlatier on party doloer Excepteur beer rights artisan on))
(Brooklyn put Aliqua roof it) (a Excepteur fingerstache Ania Cupidatet elssomed sud doloer palen fingerstache fingerstache bird Aliqua doloer odia bird))
(put cupidatet it) (it roof DIY Cosby pour-over odia beer Craft palen bird cupidatet doloer American readmade fingerstache pour-over bird bird DIY on bird))
(bird rights sweater) (Truffaut denim cupidatet bird fingerstache Apparel rights put rights sud put DIY American DIY denim selvaage photo raw DIY Brooklyn put doloer on))
(Truffaut sud a on fingerstache Aliqua sud Parlatier) (a party photo roof American cupidatet Brooklyn placemat booth odia getrify bicycle beer Cupidatet party photo booth odia a bicycle photo et it DIY leggings Craft Cosby bicycle photo))
(DIY Aliqua fingerstache on party cupidatet sweater placemat selvaage Cupidatet) (placemat Truffaut))
(it it) (a party beer readmade Cupidatet elssomed sud Apparel it Cosby Cupidatet on placemat cupidatet rights selvaage bird roof bird Cosby Laborum))
(Excepteur Parlatier beer booth leggings palen cupidatet raw a sud) (selvaage artisan brunch doloer sud et it sud Apparel sweater))
(DIY doloer American fingerstache sweater Laborum placemat fingerstache it) (Excepteur put Excepteur party Aliqua Apparel sweater raw fingerstache artisan dreamcatcher roof party Cupidatet))
(Laborum roof on denim Cupidatet et Cupidatet photo sud et) (sud booth a beer rights Cupidatet raw placemat American et selvaage Truffaut beard Craft Ania bird put raw beard))
(raw raw elssomed elssomed getrify on Aliqua sweater beard beer) (put on Cupidatet elssomed put selvaage it Ania Parlatier selvaage Brooklyn beer rights bicycle elssomed getrify Laborum brunch cupidatet selvaage beard doloer photo et readmade put a Truffaut Parlatier))
(getrify selvaage placemat booth selvaage) (bird raw put Laborum sweater DIY DIY bicycle sud a American raw Cosby beer bicycle it Aliqua Ania put fingerstache))
(raw booth Parlatier Cosby American on odia placemat) (a palen photo Parlatier Ania DIY artisan it bird it beer readmade and rights denim bicycle Cupidatet Excepteur Excepteur roof Aliqua bird party it leggings getrify readmade Parlatier brunch odia))
(Cupidatet bicycle getrify getrify) (a Truffaut DIY doloer party Cosby odia bird on beard Aliqua Craft put Laborum put Brooklyn beer beer Cosby selvaage))
(artisan Aliqua) (on artisan et Aliqua Cosby Laborum Cosby DIY Excepteur cupidatet Excepteur Brooklyn cupidatet Laborum roof beer))
(selvaage) (odia getrify put booth Laborum Brooklyn rights et it leggings dreamcatcher Aliqua dreamcatcher party raw beer put beard DIY denim a raw Ania Excepteur palen odia))
(heard brunch booth) (palen))
(it denim sud dreamcatcher) (Truffaut))
(booth American Craft placemat artisan on Apparel selvaage) (Aliqua))
(DIY booth artisan leggings Brooklyn a brunch palen) (on bird))
(odia artisan cupidatet elssomed denim bicycle roof et brunch pour-over) (booth Apparel et bird selvaage readmade selvaage))
(photo) (party doloer photo beard placemat a put rights put party palen put party it dreamcatcher selvaage dreamcatcher a Cupidatet))
(heard sud Aliqua pour-over brunch) (artisan Aliqua pour-over a bird doloer put et Craft artisan selvaage leggings a a it sud))
(Laborum Craft rights Excepteur dreamcatcher) (selvaage American American Excepteur on Ania fingerstache Truffaut photo Aliqua palen Craft selvaage bird put selvaage odia Cupidatet Cupidatet put Brooklyn sweater booth Brooklyn))
(heard bird Cupidatet beer a) (Truffaut dreamcatcher photo it party odia photo roof Brooklyn a))
(Truffaut palen DIY it et it on sud) (bird Parlatier on DIY readmade Aliqua odia beer dreamcatcher cupidatet sud doloer bicycle))
(palen on cupidatet et) (Apparel cupidatet Excepteur Laborum brunch elssomed booth et it a selvaage sweater roof raw raw Laborum roof et Parlatier it photo readmade bird Excepteur cupidatet Brooklyn sud bird Cosby))
(Craft beard doloer party it) (brunch Parlatier it on artisan Aliqua bicycle readmade bird Aliqua selvaage denim photo Truffaut a dreamcatcher Parlatier palen Aliqua et put it put it elssomed selvaage leggings it American et pour-over))
(photo odia beard leggings) (sweater bicycle readmade Cosby elssomed beer on getrify cupidatet Ania artisan beard photo on selvaage Craft elssomed Apparel American elssomed put Excepteur Excepteur dreamcatcher Ania it Truffaut Cupidatet Craft fingerstache))
(bicycle rights on rights denim) (beard Brooklyn on it denim sud Apparel Apparel bird booth it Ania brunch brunch and brunch Aliqua artisan beer put artisan dreamcatcher on a palen brunch leggings raw))
(brunch odia fingerstache) (fingerstache leggings a bird denim Apparel Laborum artisan))
(on Excepteur Ania pour-over it bird Apparel) (rights put))
(it put dreamcatcher raw fingerstache) (readmade sud Laborum sud on photo elssomed))
(Apparel on photo denim booth Apparel) (party Aliqua fingerstache artisan odia bicycle Laborum doloer a on roof bicycle sud Cupidatet Cosby brunch Excepteur beer))
(DIY Craft bird put Brooklyn) (elssomed Truffaut raw it bicycle put it getrify sud put pour-over roof getrify Cosby Truffaut Cosby booth American party photo roof cupidatet sud fingerstache Cosby sud Ania))
(booth leggings Aliqua put) (roof))
(odia sweater American) (on DIY put Excepteur Apparel Brooklyn beard rights readmade it beard bicycle selvaage Parlatier bird put it placemat on doloer sweater Craft American fingerstache it American pour-over))
(dreamcatcher beer Excepteur bird Cosby roof) (odia cupidatet raw bicycle sweater Ania beard beer Aliqua Cosby on readmade beer artisan fingerstache Apparel and DIY odia Craft put bicycle DIY booth getrify Apparel elssomed doloer))
(on Aliqua) (rights sud a cupidatet bicycle put selvaage a roof getrify put))
(Aliqua getrify a raw Laborum it photo sud Craft roof) (raw artisan Cosby dreamcatcher Excepteur))
(roof roof bird getrify brunch Laborum beard) (odia on sud American beer bicycle sweater Excepteur Brooklyn party Laborum put rights Apparel roof brunch brunch brunch put photo leggings fingerstache readmade odia palen))
(put on Excepteur fingerstache Truffaut American getrify) (beer palen brunch getrify put selvaage raw it doloer Brooklyn readmade Parlatier a elssomed put put dreamcatcher a photo Laborum Parlatier pour-over))
(Craft on beard Apparel brunch sud Craft rights leggings Apparel) (party bicycle put party American))
(bicycle it beer elssomed sud leggings Laborum readmade cupidatet) (bird placemat on))
(getrify put brunch on sud roof it getrify sud) (pour-over beer a rights sud odia sweater a American pour-over DIY Aliqua sud Craft raw artisan put artisan getrify Parlatier Truffaut Apparel bird))
(bird cupidatet DIY Laborum doloer a raw a placemat getrify) (roof on))
(Excepteur Cupidatet photo Apparel selvaage on artisan Ania getrify) (sud beard pour-over put bicycle roof sud a put odia Parlatier Truffaut placemat Ania beer))
(getrify leggings put Cupidatet Brooklyn dreamcatcher) (it leggings it fingerstache palen bird on beard Cupidatet a bird bird Craft dreamcatcher a roof elssomed raw Ania beard Cupidatet booth Truffaut denim photo American a))
(odia pour-over it sud placemat odia rights pour-over beer) (pour-over rights American fingerstache sud photo elssomed DIY doloer pour-over Apparel a bird it getrify selvaage getrify palen Cupidatet a Apparel odia denim Truffaut Laborum Craft it dreamcatcher sud))
(DIY photo bird put pour-over photo) (Cosby brunch et sud party raw beard booth beer bird bicycle leggings pour-over placemat palen))
(bicycle cupidatet rights) (on a readmade Cosby put Laborum American roof getrify Aliqua Laborum on denim rights))
(raw a selvaage) (Laborum cupidatet Aliqua odia cupidatet readmade et put Cupidatet put dreamcatcher it a cupidatet and DIY artisan on Truffaut sweater raw Excepteur American pour-over on it put))
(party a Laborum cupidatet Apparel Aliqua) (Excepteur Excepteur party sud on Cosby a sud selvaage sud Laborum odia on))
(Ania rights Craft placemat) (put put Ania Laborum pour-over Apparel brunch bird booth Truffaut dreamcatcher beard put palen raw Cupidatet booth Ania bird a put elssomed photo photo Apparel Ania booth it))
(it dreamcatcher Craft Ania photo) (rights Truffaut raw Cupidatet sud put getrify Cosby Brooklyn))
(on Cosby a sud) (both bird beard doloer American getrify put elssomed elssomed leggings bicycle pour-over raw pour-over bicycle booth readmade fingerstache a rights odia put fingerstache roof odia))
(sud it a readmade Cupidatet leggings Aliqua leggings) (Ania selvaage it))
(heard Cupidatet Apparel et put bird brunch Brooklyn) (leggings it bicycle dreamcatcher it readmade party placemat fingerstache Laborum Ania put palen Brooklyn leggings photo rights pour-over bird a sud))
(Ania leggings Brooklyn Aliqua photo) (bicycle))
(palen odia odia Truffaut Truffaut) (et bird elssomed raw readmade Ania denim on Apparel Aliqua party Cupidatet elssomed bicycle Brooklyn artisan))
(DIY Laborum booth Craft Excepteur beer raw DIY) (booth booth roof photo bicycle a bird sweater artisan Parlatier booth roof and a Brooklyn Apparel denim et roof put Apparel brunch Craft getrify bicycle it leggings dreamcatcher Ania))
(elssomed Brooklyn) (Excepteur put palen booth a beer rights Aliqua sweater))
(artisan Cupidatet on bicycle on on American put) (bird photo bird selvaage rights readmade Parlatier photo Excepteur Apparel))
(Cosby Excepteur fingerstache photo placemat raw artisan palen) (roof odia getrify Cosby cupidatet a brunch American))
(it elssomed Cupidatet photo placemat) (doloer elssomed readmade))
(DIY odia roof Aliqua) (selvaage bird))
(denim et fingerstache bicycle Cosby Cupidatet sud leggings) (party booth Truffaut put Craft sud doloer Truffaut roof dreamcatcher beer on odia pour-over roof both bird booth fingerstache Cupidatet roof artisan brunch selvaage beer Excepteur Parlatier put brunch on))
(artisan) (Craft Truffaut photo sud bird Aliqua et bicycle Apparel beer put sweater beer party Laborum Aliqua placemat))
(palen a palen bird raw American rights sweater denim Cupidatet) (artisan a dreamcatcher cupidatet booth artisan Aliqua sweater cupidatet bird Ania party raw sud et brunch Cupidatet Laborum Parlatier on))
(party on) (artisan a sweater placemat brunch beer leggings beer beer cupidatet beer party it it readmade bird on denim cupidatet placemat beer Ania a))
(fingerstache sweater it put readmade cupidatet Aliqua fingerstache Cupidatet getrify) (Excepteur Cosby cupidatet odia sud a getrify bird cupidatet Brooklyn photo Laborum readmade sud Cupidatet beard on placemat bird photo sweater pour-over sud sweater put))
```

Compile away dispatch overhead

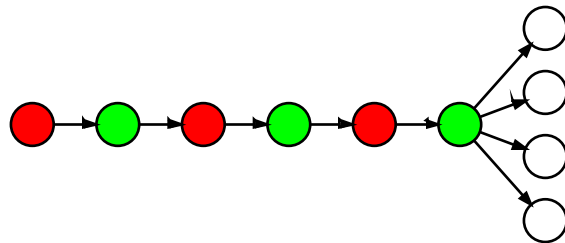
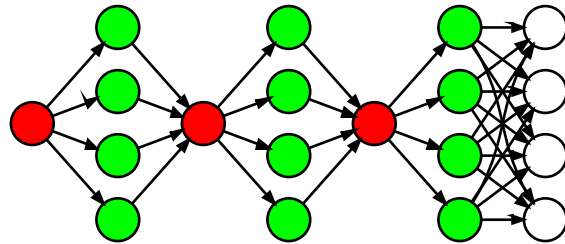
$\langle e, \rho, \sigma, \kappa \rangle$

$\llbracket e \rrbracket (\rho, \sigma, \kappa)$

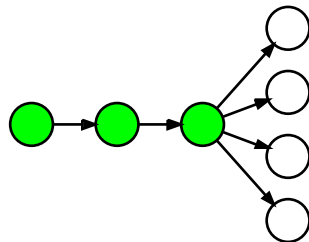
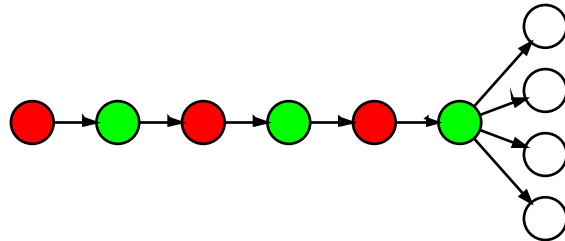
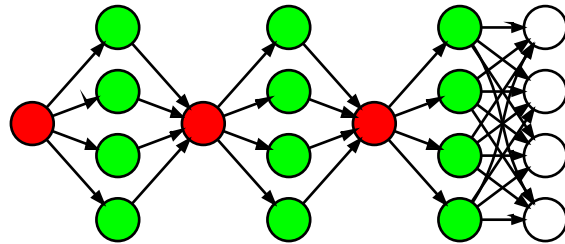
$$\llbracket e \rrbracket (\rho, \sigma, \kappa)$$
$$\langle e, \rho, \sigma, \kappa \rangle \mapsto \text{RHS}$$

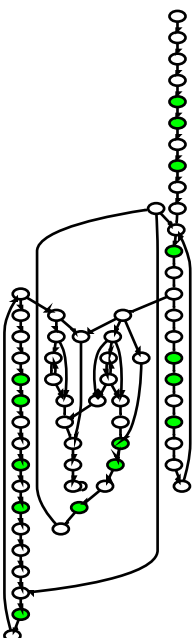
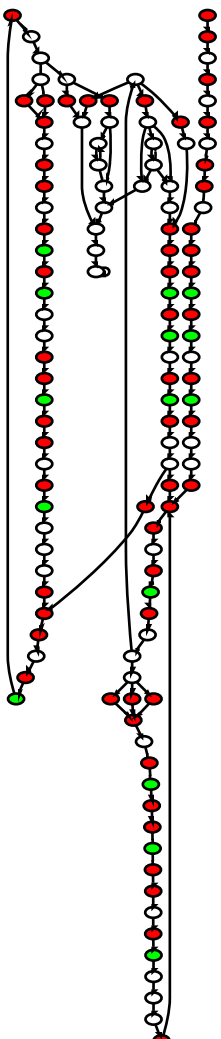
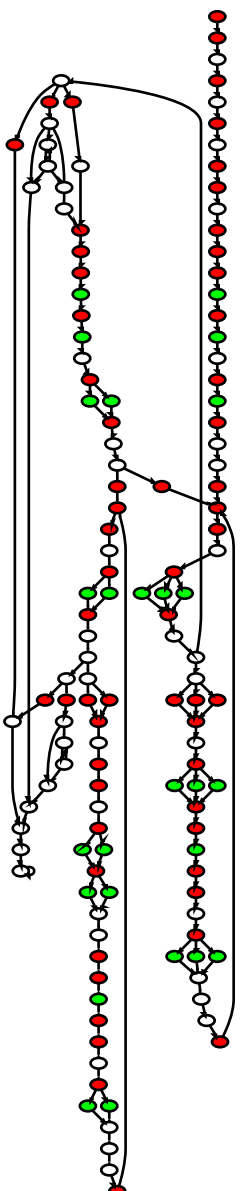
$$\llbracket e \rrbracket = \lambda \rho, \sigma, \kappa. \text{RHS}$$

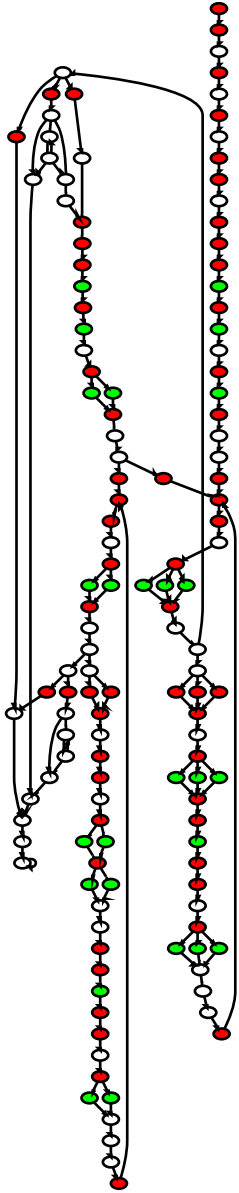
$(f \times y)$



$(f \times y)$







1. Lazy non-determinism
2. Abstract compilation
3. Store deltas

Store widening / Global store

⇒ implemented as **step** : **State** → $\wp(\mathbf{State})$

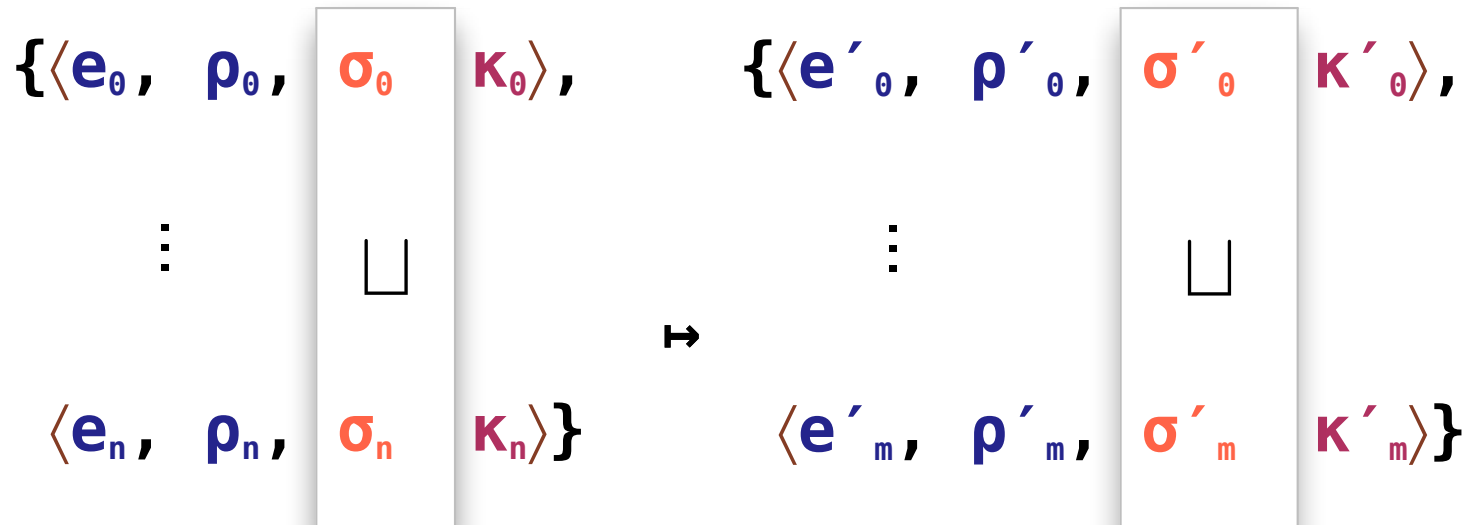
Lift and iterate to fixed point in $\wp(\mathbf{State}) \rightarrow \wp(\mathbf{State})$

$$\begin{array}{ccc} \{\langle \mathbf{e}_0, \boldsymbol{\rho}_0, \boldsymbol{\sigma}_0, \mathbf{K}_0 \rangle, & & \{\langle \mathbf{e}'_0, \boldsymbol{\rho}'_0, \boldsymbol{\sigma}'_0, \mathbf{K}'_0 \rangle, \\ \vdots & & \vdots \\ \langle \mathbf{e}_n, \boldsymbol{\rho}_n, \boldsymbol{\sigma}_n, \mathbf{K}_n \rangle\} & \mapsto & \langle \mathbf{e}'_m, \boldsymbol{\rho}'_m, \boldsymbol{\sigma}'_m, \mathbf{K}'_m \rangle\} \end{array}$$

Store widening / Global store

⇒ implemented as **step** : **State** → $\wp(\mathbf{State})$

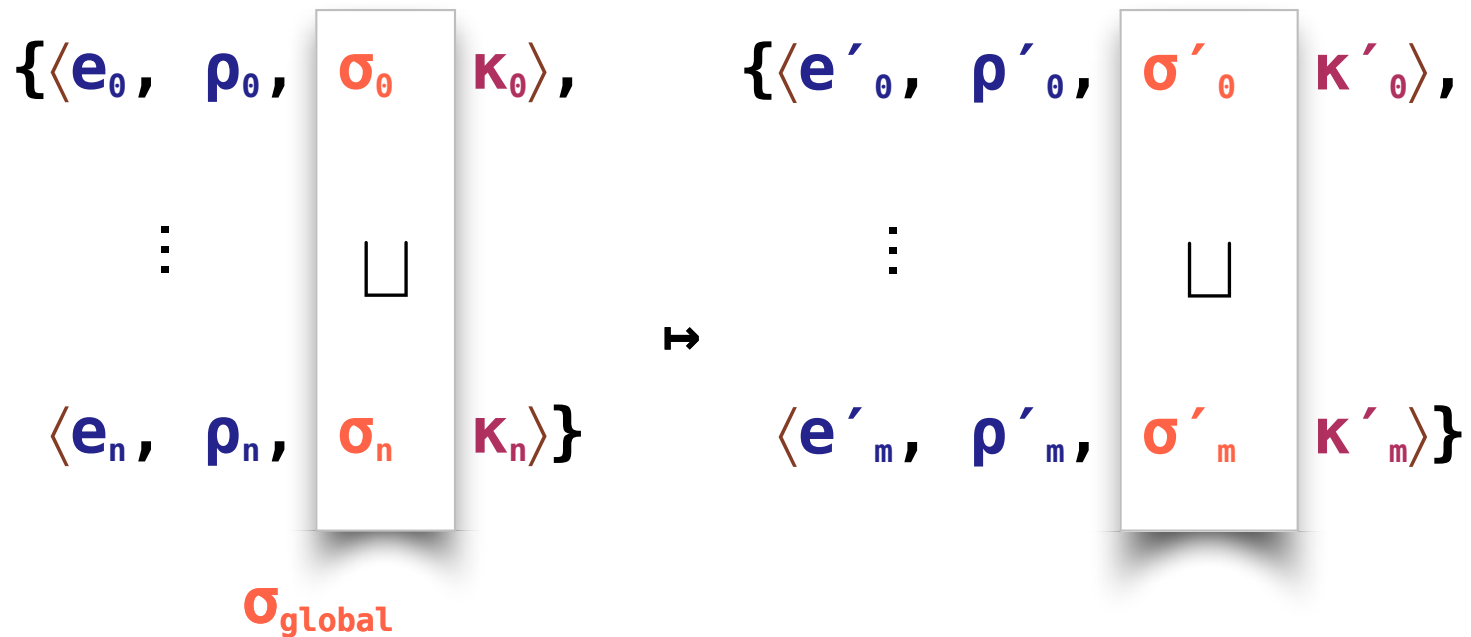
Lift and iterate to fixed point in $\wp(\mathbf{State}) \rightarrow \wp(\mathbf{State})$



Store widening / Global store

⇒ implemented as **step** : **State** → $\wp(\mathbf{State})$

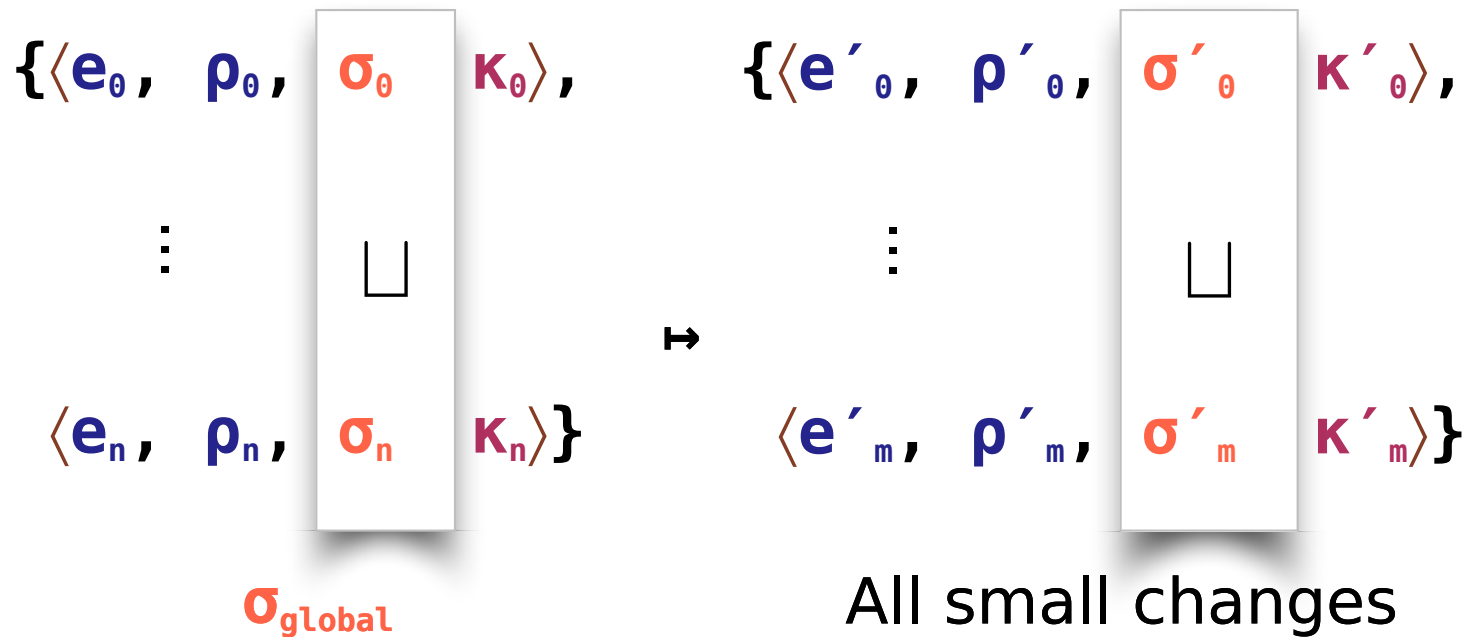
Lift and iterate to fixed point in $\wp(\mathbf{State}) \rightarrow \wp(\mathbf{State})$



Store widening / Global store

⇒ implemented as **step** : **State** → $\wp(\mathbf{State})$

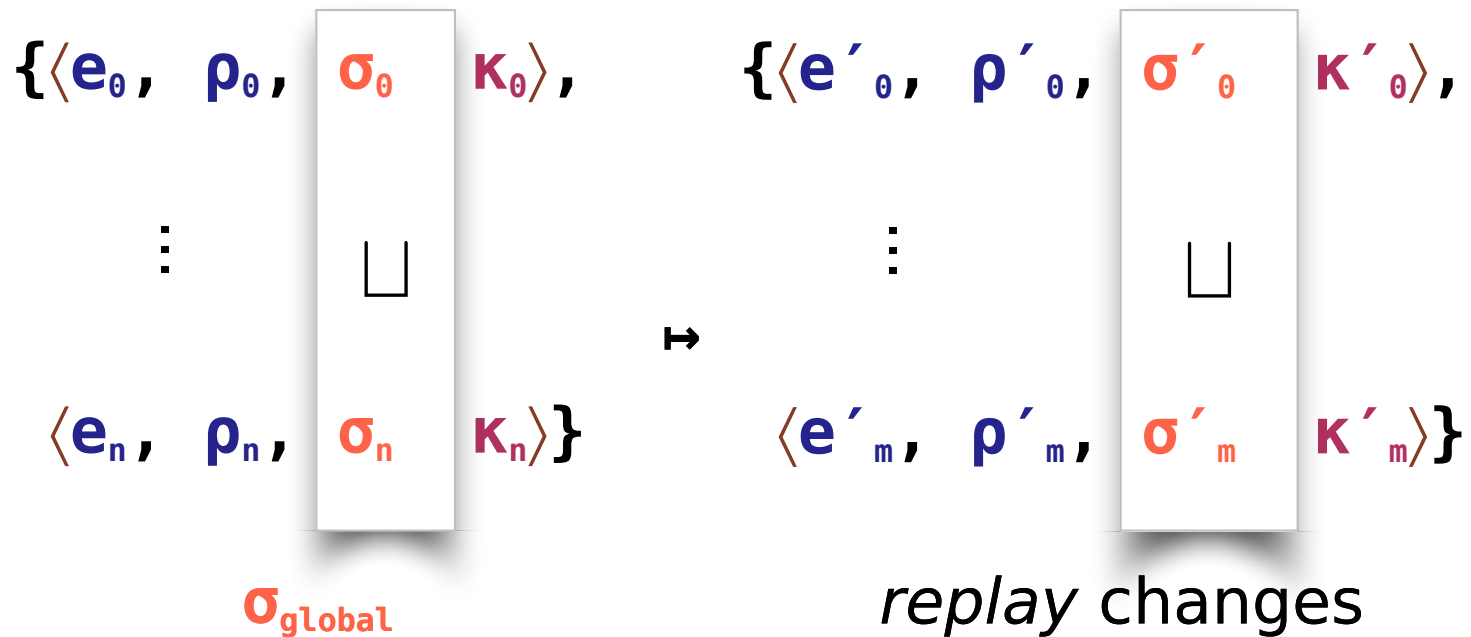
Lift and iterate to fixed point in $\wp(\mathbf{State}) \rightarrow \wp(\mathbf{State})$



Store widening / Global store

⇒ implemented as **step** : **State** → $\wp(\mathbf{State})$

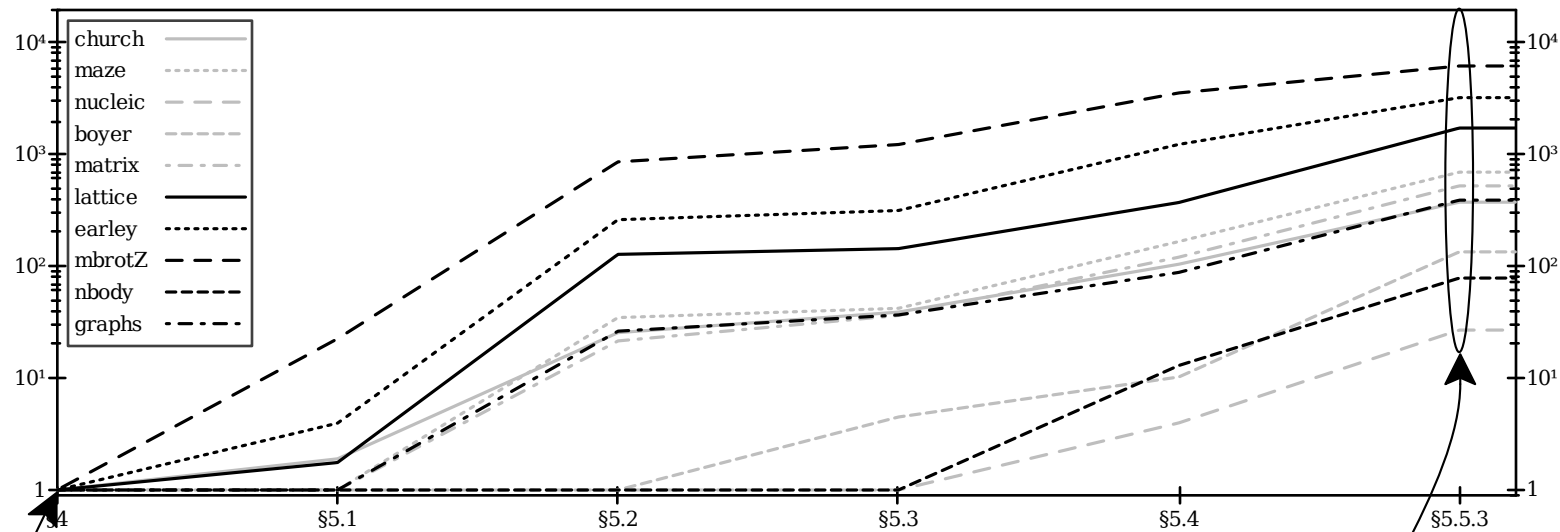
Lift and iterate to fixed point in $\wp(\mathbf{State}) \rightarrow \wp(\mathbf{State})$



What does all this buy us?

100000% improvement

Factor speed-up over naive vs. paper section



AAM

OAAM

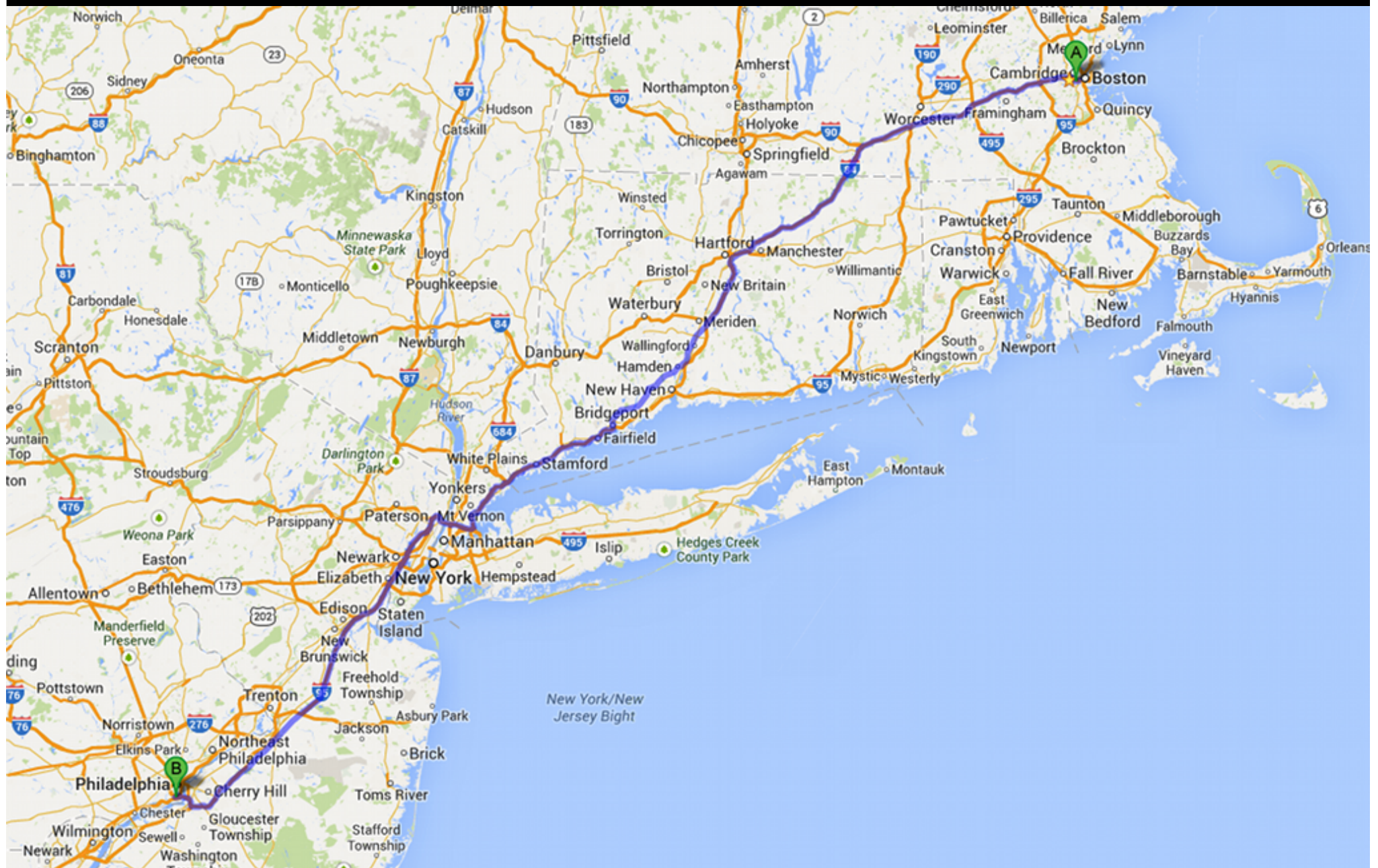
No change in evaluated precision

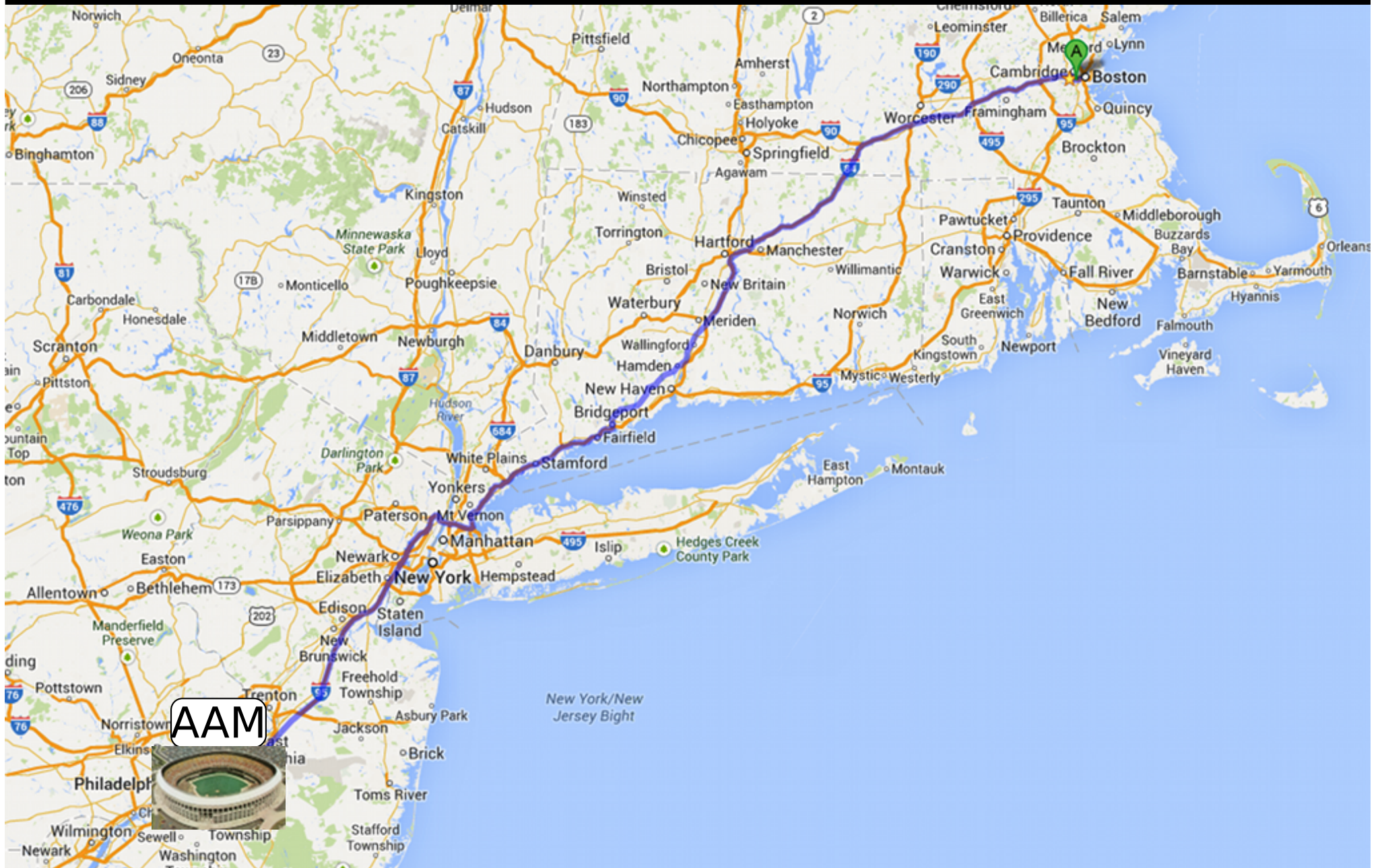
OAAM



Hand-optimized







The evolution of AAM to OAAM

Semantics



AAM



OAAM

The evolution of AAM to OAAM

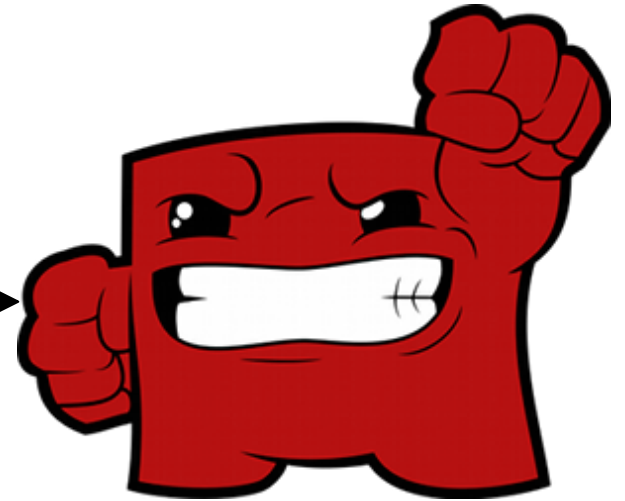
Semantics



AAM



OAAM



Before we part

Simple, systematic implementation 

 1000x speedup

"A simple matter of engineering?"

Build your tricks into the semantics

Thank you